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# Several Classes of BCI-algebras and their Properties 

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#### Abstract

Summary. I have formalized the BCI-algebras closely following the book [6], sections 1.1 to $1.3,1.6,2.1$ to 2.3 , and 2.7. In this article the general theory of BCI-algebras and several classes of BCI-algebras are given.


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The articles [10], [4], [13], [9], [3], [12], [2], [11], [5], [7], [8], [1], and [14] provide the notation and terminology for this paper.

## 1. The Basics of General Theory of BCI-algebras

We introduce BCI structures which are extensions of 1-sorted structure and are systems
$\langle$ a carrier, an internal complement $\rangle$,
where the carrier is a set and the internal complement is a binary operation on the carrier.

Let us note that there exists a BCI structure which is non empty and strict.
Let $A$ be a BCI structure and let $x, y$ be elements of $A$. The functor $x \backslash y$ yielding an element of $A$ is defined by:
(Def. 1) $x \backslash y=($ the internal complement of $A)(x, y)$.
We introduce BCI structures with 0 which are extensions of BCI structure and zero structure and are systems
$\langle$ a carrier, an internal complement, a zero 〉,
where the carrier is a set, the internal complement is a binary operation on the carrier, and the zero is an element of the carrier.

Let us note that there exists a BCI structure with 0 which is non empty and strict.

Let $I_{1}$ be a non empty BCI structure with 0 and let $x$ be an element of $I_{1}$. The functor $x^{\mathrm{c}}$ yields an element of $I_{1}$ and is defined by:
(Def. 2) $\quad x^{\mathrm{c}}=0_{\left(I_{1}\right)} \backslash x$.
Let $I_{1}$ be a non empty BCI structure with 0 . We say that $I_{1}$ is B if and only if:
(Def. 3) For all elements $x, y, z$ of $I_{1}$ holds $x \backslash y \backslash(z \backslash y) \backslash(x \backslash z)=0_{\left(I_{1}\right)}$.
We say that $I_{1}$ is C if and only if:
(Def. 4) For all elements $x, y, z$ of $I_{1}$ holds $x \backslash y \backslash z \backslash(x \backslash z \backslash y)=0_{\left(I_{1}\right)}$.
We say that $I_{1}$ is I if and only if:
(Def. 5) For every element $x$ of $I_{1}$ holds $x \backslash x=0_{\left(I_{1}\right)}$.
We say that $I_{1}$ is K if and only if:
(Def. 6) For all elements $x, y$ of $I_{1}$ holds $x \backslash y \backslash x=0_{\left(I_{1}\right)}$.
We say that $I_{1}$ is BCI-4 if and only if:
(Def. 7) For all elements $x, y$ of $I_{1}$ such that $x \backslash y=0_{\left(I_{1}\right)}$ and $y \backslash x=0_{\left(I_{1}\right)}$ holds $x=y$.
We say that $I_{1}$ is BCK-5 if and only if:
(Def. 8) For every element $x$ of $I_{1}$ holds $x^{\mathrm{c}}=0_{\left(I_{1}\right)}$.
The BCI structure BCI-EXAMPLE with 0 is defined as follows:
(Def. 9) BCI-EXAMPLE $=\left\langle\{\emptyset\}, \mathrm{op}_{2}, \mathrm{op}_{0}\right\rangle$.
Let us note that BCI-EXAMPLE is strict and non empty.
One can verify that there exists a non empty BCI structure with 0 which is strict, B, C, I, and BCI-4.

A BCI-algebra is B C I BCI-4 non empty BCI structure with 0.
Let $X$ be a BCI-algebra. A BCI-algebra is called a subalgebra of $X$ if it satisfies the conditions (Def. 10).
(Def. 10)(i) $\quad 0_{i t}=0_{X}$,
(ii) the carrier of it $\subseteq$ the carrier of $X$, and
(iii) the internal complement of it $=($ the internal complement of $X) \upharpoonright($ the carrier of it).
The following proposition is true
(1) Let $X$ be a non empty BCI structure with 0 . Then $X$ is a BCI-algebra if and only if the following conditions are satisfied:
(i) $X$ is I and BCI-4, and
(ii) for all elements $x, y, z$ of $X$ holds $x \backslash y \backslash(x \backslash z) \backslash(z \backslash y)=0_{X}$ and $x \backslash(x \backslash y) \backslash y=0_{X}$.

One can check that there exists a BCI-algebra which is strict and BCK-5.
A BCK-algebra is BCK-5 BCI-algebra.
Let $I_{1}$ be a non empty BCI structure with 0 and let $x, y$ be elements of $I_{1}$. The predicate $x \leq y$ is defined as follows:
(Def. 11) $x \backslash y=0_{\left(I_{1}\right)}$.
We use the following convention: $X$ denotes a BCI-algebra, $x, y, z, u, a, b$ denote elements of $X$, and $I_{1}$ denotes a non empty subset of $X$.

We now state a number of propositions:
(2) $x \backslash 0_{X}=x$.
(3) If $x \backslash y=0_{X}$ and $y \backslash z=0_{X}$, then $x \backslash z=0_{X}$.
(4) If $x \backslash y=0_{X}$, then $x \backslash z \backslash(y \backslash z)=0_{X}$ and $z \backslash y \backslash(z \backslash x)=0_{X}$.
(5) If $x \leq y$, then $x \backslash z \leq y \backslash z$ and $z \backslash y \leq z \backslash x$.
(6) If $x \backslash y=0_{X}$, then $(y \backslash x)^{\mathrm{c}}=0_{X}$.
(7) $x \backslash y \backslash z=x \backslash z \backslash y$.
(8) $x \backslash(x \backslash(x \backslash y))=x \backslash y$.
(9) $\quad(x \backslash y)^{\mathrm{c}}=x^{\mathrm{c}} \backslash y^{\mathrm{c}}$.
(10) $x \backslash(x \backslash y) \backslash(y \backslash x) \backslash(x \backslash(x \backslash(y \backslash(y \backslash x))))=0_{X}$.
(11) Let $X$ be a non empty BCI structure with 0 . Then $X$ is a BCI-algebra if and only if the following conditions are satisfied:
(i) $X$ is BCI-4, and
(ii) for all elements $x, y, z$ of $X$ holds $x \backslash y \backslash(x \backslash z) \backslash(z \backslash y)=0_{X}$ and $x \backslash 0_{X}=x$.
(12) If for every BCI-algebra $X$ and for all elements $x, y$ of $X$ holds $x \backslash(x \backslash y)=$ $y \backslash(y \backslash x)$, then $X$ is a BCK-algebra.
(13) If for every BCI-algebra $X$ and for all elements $x, y$ of $X$ holds $x \backslash y \backslash y=$ $x \backslash y$, then $X$ is a BCK-algebra.
(14) If for every BCI-algebra $X$ and for all elements $x, y$ of $X$ holds $x \backslash(y \backslash x)=$ $x$, then $X$ is a BCK-algebra.
(15) If for every BCI-algebra $X$ and for all elements $x, y, z$ of $X$ holds ( $x$ \} $y) \backslash y=x \backslash z \backslash(y \backslash z)$, then $X$ is a BCK-algebra.
(16) If for every BCI-algebra $X$ and for all elements $x, y$ of $X$ holds $x \backslash y \backslash$ $(y \backslash x)=x \backslash y$, then $X$ is a BCK-algebra.
(17) If for every BCI-algebra $X$ and for all elements $x, y$ of $X$ holds $x \backslash y \backslash$ $(x \backslash y \backslash(y \backslash x))=0_{X}$, then $X$ is a BCK-algebra.
(18) For every BCI-algebra $X$ holds $X$ is K iff $X$ is a BCK-algebra.

Let $X$ be a BCI-algebra. The functor BCK-part $X$ yielding a non empty subset of $X$ is defined by:
(Def. 12) BCK-part $X=\left\{x ; x\right.$ ranges over elements of $\left.X: 0_{X} \leq x\right\}$.

Next we state the proposition
(19) $0_{X} \in$ BCK-part $X$.

Let us consider $X$. Note that $0_{X}$
Next we state three propositions:
(20) For all elements $x, y$ of BCK-part $X$ holds $x \backslash y \in$ BCK-part $X$.
(21) For every element $x$ of $X$ and for every element $y$ of BCK-part $X$ holds $x \backslash y \leq x$.
(22) $X$ is a subalgebra of $X$.

Let $X$ be a BCI-algebra and let $I_{1}$ be a subalgebra of $X$. We say that $I_{1}$ is proper if and only if:
(Def. 13) $\quad I_{1} \neq X$.
Let us consider $X$. Note that there exists a subalgebra of $X$ which is non proper.

Let $X$ be a BCI-algebra and let $I_{1}$ be an element of $X$. We say that $I_{1}$ is atom if and only if:
(Def. 14) For every element $z$ of $X$ such that $z \backslash I_{1}=0_{X}$ holds $z=I_{1}$.
Let $X$ be a BCI-algebra. The functor AtomSet $X$ yields a non empty subset of $X$ and is defined by:
(Def. 15) AtomSet $X=\{x ; x$ ranges over elements of $X: x$ is atom $\}$.
One can prove the following propositions:
(23) $0_{X} \in$ AtomSet $X$.
(24) For every element $x$ of $X$ holds $x \in$ AtomSet $X$ iff for every element $z$ of $X$ holds $z \backslash(z \backslash x)=x$.
(25) For every element $x$ of $X$ holds $x \in$ AtomSet $X$ iff for all elements $u, z$ of $X$ holds $z \backslash u \backslash(z \backslash x)=x \backslash u$.
(26) For every element $x$ of $X$ holds $x \in$ AtomSet $X$ iff for all elements $y, z$ of $X$ holds $x \backslash(z \backslash y) \leq y \backslash(z \backslash x)$.
(27) For every element $x$ of $X$ holds $x \in$ AtomSet $X$ iff for all elements $y, z$, $u$ of $X$ holds $(x \backslash u) \backslash(z \backslash y) \leq y \backslash u \backslash(z \backslash x)$.
(28) For every element $x$ of $X$ holds $x \in$ AtomSet $X$ iff for every element $z$ of $X$ holds $z^{\mathrm{c}} \backslash x^{\mathrm{c}}=x \backslash z$.
(29) For every element $x$ of $X$ holds $x \in$ AtomSet $X$ iff $\left(x^{\mathrm{c}}\right)^{\mathrm{c}}=x$.
(30) For every element $x$ of $X$ holds $x \in$ AtomSet $X$ iff for every element $z$ of $X$ holds $(z \backslash x)^{\mathrm{c}}=x \backslash z$.
(31) For every element $x$ of $X$ holds $x \in$ AtomSet $X$ iff for every element $z$ of $X$ holds $\left((x \backslash z)^{\mathrm{c}}\right)^{\mathrm{c}}=x \backslash z$.
(32) For every element $x$ of $X$ holds $x \in$ AtomSet $X$ iff for all elements $z, u$ of $X$ holds $z \backslash(z \backslash(x \backslash u))=x \backslash u$.
(33) For every element $a$ of AtomSet $X$ and for every element $x$ of $X$ holds $a \backslash x \in$ AtomSet $X$.
Let $X$ be a BCI-algebra and let $a, b$ be elements of AtomSet $X$. Then $a \backslash b$ is an element of AtomSet $X$.

One can prove the following propositions:
(34) For every element $x$ of $X$ holds $x^{\mathrm{c}} \in$ AtomSet $X$.
(35) For every element $x$ of $X$ there exists an element $a$ of AtomSet $X$ such that $a \leq x$.
Let $X$ be a BCI-algebra. We say that $X$ is generated by atom if and only if:
(Def. 16) For every element $x$ of $X$ there exists an element $a$ of AtomSet $X$ such that $a \leq x$.
Let $X$ be a BCI-algebra and let $a$ be an element of AtomSet $X$. The functor BranchV $a$ yields a non empty subset of $X$ and is defined as follows:
(Def. 17) BranchV $a=\{x ; x$ ranges over elements of $X: a \leq x\}$.
We now state several propositions:
(36) Every BCI-algebra is generated by atom.
(37) For all elements $a, b$ of AtomSet $X$ and for every element $x$ of BranchV $b$ holds $a \backslash x=a \backslash b$.
(38) For every element $a$ of AtomSet $X$ and for every element $x$ of BCK-part $X$ holds $a \backslash x=a$.
(39) For all elements $a, b$ of AtomSet $X$ and for every element $x$ of $\operatorname{BranchV} a$ and for every element $y$ of $\operatorname{BranchV} b$ holds $x \backslash y \in \operatorname{BranchV}(a \backslash b)$.
(40) For every element $a$ of AtomSet $X$ and for all elements $x, y$ of BranchV $a$ holds $x \backslash y \in$ BCK-part $X$.
(41) For all elements $a, b$ of AtomSet $X$ and for every element $x$ of BranchV $a$ and for every element $y$ of BranchV $b$ such that $a \neq b$ holds $x \backslash y \notin$ BCK-part $X$.
(42) For all elements $a, b$ of AtomSet $X$ such that $a \neq b$ holds BranchV $a \cap$ BranchV $b=\emptyset$.
Let $X$ be a BCI-algebra. A non empty subset of $X$ is said to be an ideal of $X$ if:
(Def. 18) $0_{X} \in$ it and for all elements $x, y$ of $X$ such that $x \backslash y \in$ it and $y \in$ it holds $x \in$ it.
Let $X$ be a BCI-algebra and let $I_{1}$ be an ideal of $X$. We say that $I_{1}$ is closed if and only if:
(Def. 19) For every element $x$ of $I_{1}$ holds $x^{\mathrm{c}} \in I_{1}$.
Let us consider $X$. Note that there exists an ideal of $X$ which is closed.
Next we state four propositions:
(43) $\left\{0_{X}\right\}$ is a closed ideal of $X$.
(44) The carrier of $X$ is a closed ideal of $X$.
(45) BCK-part $X$ is a closed ideal of $X$.
(46) If $I_{1}$ is an ideal of $X$, then for all elements $x, y$ of $X$ such that $x \in I_{1}$ and $y \leq x$ holds $y \in I_{1}$.

## 2. Associative BCI-Algebras

Let $I_{1}$ be a BCI-algebra. We say that $I_{1}$ is associative if and only if:
(Def. 20) For all elements $x, y, z$ of $I_{1}$ holds $(x \backslash y) \backslash z=x \backslash(y \backslash z)$.
We say that $I_{1}$ is quasi-associative if and only if:
(Def. 21) For every element $x$ of $I_{1}$ holds $\left(x^{\mathrm{c}}\right)^{\mathrm{c}}=x^{\mathrm{c}}$.
We say that $I_{1}$ is positive-implicative if and only if:
(Def. 22) For all elements $x, y$ of $I_{1}$ holds $(x \backslash(x \backslash y)) \backslash(y \backslash x)=x \backslash(x \backslash(y \backslash(y \backslash x)))$.
We say that $I_{1}$ is weakly-positive-implicative if and only if:
(Def. 23) For all elements $x, y, z$ of $I_{1}$ holds $(x \backslash y) \backslash z=x \backslash z \backslash z \backslash(y \backslash z)$.
We say that $I_{1}$ is implicative if and only if:
(Def. 24) For all elements $x, y$ of $I_{1}$ holds $(x \backslash(x \backslash y)) \backslash(y \backslash x)=y \backslash(y \backslash x)$.
We say that $I_{1}$ is weakly-implicative if and only if:
(Def. 25) For all elements $x, y$ of $I_{1}$ holds $x \backslash(y \backslash x) \backslash(y \backslash x)^{\mathrm{c}}=x$.
We say that $I_{1}$ is $p$-semisimple if and only if:
(Def. 26) For all elements $x, y$ of $I_{1}$ holds $x \backslash(x \backslash y)=y$.
We say that $I_{1}$ is alternative if and only if:
(Def. 27) For all elements $x, y$ of $I_{1}$ holds $x \backslash(x \backslash y)=(x \backslash x) \backslash y$ and $(x \backslash y) \backslash y=$ $x \backslash(y \backslash y)$.
One can check that there exists a BCI-algebra which is implicative, positiveimplicative, $p$-semisimple, associative, weakly-implicative, and weakly-positiveimplicative.

Next we state several propositions:
(47) $X$ is associative iff for every element $x$ of $X$ holds $x^{\mathrm{c}}=x$.
(48) For all elements $x, y$ of $X$ holds $y \backslash x=x \backslash y$ iff $X$ is associative.
(49) Let $X$ be a non empty BCI structure with 0 . Then $X$ is an associative BCI-algebra if and only if for all elements $x, y, z$ of $X$ holds $y \backslash x \backslash(z \backslash x)=$ $z \backslash y$ and $x \backslash 0_{X}=x$.
(50) Let $X$ be a non empty BCI structure with 0 . Then $X$ is an associative BCI-algebra if and only if for all elements $x, y, z$ of $X$ holds $x \backslash y \backslash(x \backslash z)=$ $z \backslash y$ and $x^{\mathrm{c}}=x$.
(51) Let $X$ be a non empty BCI structure with 0 . Then $X$ is an associative BCI-algebra if and only if for all elements $x, y, z$ of $X$ holds $x \backslash y \backslash(x \backslash z)=$ $y \backslash z$ and $x \backslash 0_{X}=x$.

## 3. $p$-SEMISIMPLE BCI-ALGEBRAS

One can prove the following propositions:
(52) $\quad X$ is $p$-semisimple iff every element of $X$ is atom.
(53) If $X$ is $p$-semisimple, then BCK-part $X=\left\{0_{X}\right\}$.
(54) $\quad X$ is $p$-semisimple iff for every element $x$ of $X$ holds $\left(x^{\mathrm{c}}\right)^{\mathrm{c}}=x$.
(55) $\quad X$ is $p$-semisimple iff for all $x, y$ holds $y \backslash(y \backslash x)=x$.
(56) $\quad X$ is $p$-semisimple iff for all $x, y, z$ holds $z \backslash y \backslash(z \backslash x)=x \backslash y$.
(57) $\quad X$ is $p$-semisimple iff for all $x, y, z$ holds $x \backslash(z \backslash y)=y \backslash(z \backslash x)$.
(58) $\quad X$ is $p$-semisimple iff for all $x, y, z, u$ holds $(x \backslash u) \backslash(z \backslash y)=y \backslash u \backslash(z \backslash x)$.
(59) $\quad X$ is $p$-semisimple iff for all $x, z$ holds $z^{\mathrm{c}} \backslash x^{\mathrm{c}}=x \backslash z$.
(60) $X$ is $p$-semisimple iff for all $x, z$ holds $\left((x \backslash z)^{\mathrm{c}}\right)^{\mathrm{c}}=x \backslash z$.
(61) $X$ is $p$-semisimple iff for all $x, u$, $z$ holds $z \backslash(z \backslash(x \backslash u))=x \backslash u$.
(62) $\quad X$ is $p$-semisimple iff for every $x$ such that $x^{\mathrm{c}}=0_{X}$ holds $x=0_{X}$.
(63) $\quad X$ is $p$-semisimple iff for all $x, y$ holds $x \backslash y^{\mathrm{c}}=y \backslash x^{\mathrm{c}}$.
(64) $\quad X$ is $p$-semisimple iff for all $x, y, z, u$ holds $(x \backslash y) \backslash(z \backslash u)=x \backslash z \backslash(y \backslash u)$.
(65) $\quad X$ is $p$-semisimple iff for all $x, y, z$ holds $x \backslash y \backslash(z \backslash y)=x \backslash z$.
(66) $X$ is $p$-semisimple iff for all $x, y, z$ holds $x \backslash(y \backslash z)=(z \backslash y) \backslash x^{\mathrm{c}}$.
(67) $\quad X$ is $p$-semisimple iff for all $x, y, z$ such that $y \backslash x=z \backslash x$ holds $y=z$.
(68) $\quad X$ is $p$-semisimple iff for all $x, y, z$ such that $x \backslash y=x \backslash z$ holds $y=z$.
(69) Let $X$ be a non empty BCI structure with 0 . Then $X$ is a $p$-semisimple BCI-algebra if and only if for all elements $x, y, z$ of $X$ holds $x \backslash y \backslash(x \backslash z)=$ $z \backslash y$ and $x \backslash 0_{X}=x$.
(70) Let $X$ be a non empty BCI structure with 0 . Then $X$ is a $p$-semisimple BCI-algebra if and only if $X$ is I and for all elements $x, y, z$ of $X$ holds $x \backslash(y \backslash z)=z \backslash(y \backslash x)$ and $x \backslash 0_{X}=x$.

## 4. Quasi-Associative BCI-ALGEBRAS

Next we state several propositions:
(71) $X$ is quasi-associative iff for every element $x$ of $X$ holds $x^{\mathrm{c}} \leq x$.
(72) $\quad X$ is quasi-associative iff for all elements $x, y$ of $X$ holds $(x \backslash y)^{\mathrm{c}}=(y \backslash x)^{\mathrm{c}}$.
(73) $\quad X$ is quasi-associative iff for all elements $x, y$ of $X$ holds $x^{\mathrm{c}} \backslash y=(x \backslash y)^{\mathrm{c}}$.
(74) $\quad X$ is quasi-associative iff for all elements $x, y$ of $X$ holds $x \backslash y \backslash(y \backslash x) \in$ BCK-part $X$.
(75) $\quad X$ is quasi-associative iff for all elements $x, y, z$ of $X$ holds $(x \backslash y) \backslash z \leq$ $x \backslash(y \backslash z)$.

## 5. Alternative BCI-Algebras

We now state several propositions:
(76) If $X$ is alternative, then $x^{\mathrm{c}}=x$ and $x \backslash(x \backslash y)=y$ and $x \backslash y \backslash y=x$.
(77) If $X$ is alternative and $x \backslash a=x \backslash b$, then $a=b$.
(78) If $X$ is alternative and $a \backslash x=b \backslash x$, then $a=b$.
(79) If $X$ is alternative and $x \backslash y=0_{X}$, then $x=y$.
(80) If $X$ is alternative and $x \backslash a \backslash b=0_{X}$, then $a=x \backslash b$ and $b=x \backslash a$.

One can check the following observations:

* every BCI-algebra which is alternative is also associative,
* every BCI-algebra which is associative is also alternative, and
* every BCI-algebra which is alternative is also implicative.

The following two propositions are true:
(81) If $X$ is alternative, then $x \backslash(x \backslash y) \backslash(y \backslash x)=x$.
(82) If $X$ is alternative, then $y \backslash(y \backslash(x \backslash(x \backslash y)))=y$.

## 6. Implicative, Positive-Implicative, and Weakly-Positive-Implicative BCI-Algebras

Let us observe that every BCI-algebra which is associative is also weakly-positive-implicative and every BCI-algebra which is $p$-semisimple is also weakly-positive-implicative.

We now state two propositions:
(83) Let $X$ be a non empty BCI structure with 0 . Then $X$ is an implicative BCI-algebra if and only if for all elements $x, y, z$ of $X$ holds $x \backslash y \backslash(x \backslash$ $z) \backslash(z \backslash y)=0_{X}$ and $x \backslash 0_{X}=x$ and $(x \backslash(x \backslash y)) \backslash(y \backslash x)=y \backslash(y \backslash x)$.
(84) $X$ is weakly-positive-implicative iff for all elements $x, y$ of $X$ holds $x \backslash y=$ $x \backslash y \backslash y \backslash y^{\mathrm{c}}$.
One can verify that every BCI-algebra which is positive-implicative is also weakly-positive-implicative and every BCI-algebra which is alternative is also weakly-positive-implicative.

One can prove the following two propositions:
(85) Suppose $X$ is a weakly-positive-implicative BCI-algebra. Let $x, y$ be elements of $X$. Then $(x \backslash(x \backslash y)) \backslash(y \backslash x)=y \backslash(y \backslash x) \backslash(y \backslash x) \backslash(x \backslash y)$.
(86) Let $X$ be a non empty BCI structure with 0 . Then $X$ is a positiveimplicative BCI-algebra if and only if for all elements $x, y, z$ of $X$ holds $x \backslash y \backslash(x \backslash z) \backslash(z \backslash y)=0_{X}$ and $x \backslash 0_{X}=x$ and $x \backslash y=x \backslash y \backslash y \backslash y^{\mathrm{c}}$ and $(x \backslash(x \backslash y)) \backslash(y \backslash x)=y \backslash(y \backslash x) \backslash(y \backslash x) \backslash(x \backslash y)$.

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# Formal Languages - Concatenation and Closure 

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#### Abstract

Summary. Formal languages are introduced as subsets of the set of all 0 -based finite sequences over a given set (the alphabet). Concatenation, the $n$-th power and closure are defined and their properties are shown. Finally, it is shown that the closure of the alphabet (understood here as the language of words of length 1) equals to the set of all words over that alphabet, and that the alphabet is the minimal set with this property. Notation and terminology were taken from [5] and [13].


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The terminology and notation used here are introduced in the following articles: [10], [4], [11], [8], [9], [2], [14], [3], [1], [6], [12], and [7].

## 1. Preliminaries

For simplicity, we follow the rules: $E$ is a set, $x$ is a set, $A, B, C, D$ are subsets of $E^{\omega}, a, b, c$ are elements of $E^{\omega}, e$ is an element of $E, i, n, n_{1}, n_{2}, m$ are natural numbers, and $p, q, r_{1}, r_{2}$ are real numbers.

Let us consider $E, a, b$. Then $a^{\wedge} b$ is an element of $E^{\omega}$.
Let us consider $E$. Then $\left\rangle_{E}\right.$ is an element of $E^{\omega}$.
Let $E$ be a non empty set and let $e$ be an element of $E$. Then $\langle e\rangle$ is an element of $E^{\omega}$.

Let us consider $E, a$. Then $\{a\}$ is a subset of $E^{\omega}$.
Let us consider $E$, let $f$ be a function from $\mathbb{N}$ into $2^{E^{\omega}}$, and let us consider $n$. Then $f(n)$ is a subset of $E^{\omega}$.

One can prove the following propositions:
(1) If $\{x\} \nsubseteq X$, then $\{x\}$ misses $X$.
(2) If $n_{1}>1$ or $n_{2}>1$, then $n_{1}+n_{2}>1$.
(3) $n>0$ iff $n \geq 1$.
(4) If $r_{1}+p \leq r_{2}+q$ and $p \geq q$, then $r_{1} \leq r_{2}$.
(5) If $n_{1}+n \leq n_{2}+1$ and $n>0$, then $n_{1} \leq n_{2}$.
(6) $n_{1}+n_{2}=1$ iff $n_{1}=1$ and $n_{2}=0$ or $n_{1}=0$ and $n_{2}=1$.
(7) $\quad a^{\frown} b=\langle x\rangle$ iff $a=\langle \rangle_{E}$ and $b=\langle x\rangle$ or $b=\langle \rangle_{E}$ and $a=\langle x\rangle$.
(8) For all finite 0-sequences $p, q$ such that $a=p^{\wedge} q$ holds $p$ is an element of $E^{\omega}$ and $q$ is an element of $E^{\omega}$.
(9) If $\langle x\rangle$ is an element of $E^{\omega}$, then $x \in E$.
(10) If len $b=n+1$, then there exist $c$, $e$ such that len $c=n$ and $b=c^{\frown}\langle e\rangle$.
(11) If $a^{\wedge} a=a$, then $a=\emptyset$.

## 2. Concatenation of Languages

Let us consider $E, A, B$. The functor $A \frown B$ yields a subset of $E^{\omega}$ and is defined by:
(Def. 1) $\quad x \in A \frown B$ iff there exist $a, b$ such that $a \in A$ and $b \in B$ and $x=a^{\frown} b$.
The following propositions are true:
(12) $A \frown B=\emptyset$ iff $A=\emptyset$ or $B=\emptyset$.
(13) $\quad A \frown\left\{\left\rangle_{E}\right\}=A\right.$ and $\left\{\left\rangle_{E}\right\} \frown A=A\right.$.
(14) $\quad A \frown B=\left\{\langle \rangle_{E}\right\}$ iff $A=\left\{\langle \rangle_{E}\right\}$ and $B=\left\{\langle \rangle_{E}\right\}$.
(15) $\left\rangle_{E} \in A \frown B\right.$ iff $\left\rangle_{E} \in A\right.$ and $\left\rangle_{E} \in B\right.$.
(16) If $\left\rangle_{E} \in B\right.$, then $A \subseteq A \frown B$ and $A \subseteq B \frown A$.
(17) If $A \subseteq C$ and $B \subseteq D$, then $A \frown B \subseteq C \frown D$.
(18) $(A \frown B) \frown C=A \frown(B \frown C)$.
(19) $\quad A \frown(B \cap C) \subseteq(A \frown B) \cap(A \frown C)$ and $(B \cap C) \frown A \subseteq(B \frown A) \cap(C \frown A)$.
(20) $\quad A \frown B \cup A \frown C=A \frown(B \cup C)$ and $B \frown A \cup C \frown A=(B \cup C) \frown A$.
(21) $A \frown B \backslash A \frown C \subseteq A \frown(B \backslash C)$ and $B \frown A \backslash C \frown A \subseteq(B \backslash C) \frown A$.
(22) $A \frown B \doteq A \frown C \subseteq A \frown(B \doteq C)$ and $B \frown A \doteq C \frown A \subseteq(B \subset C) \frown A$.

## 3. $n$-Th Power of a Language

Let us consider $E, A, n$. The functor $A^{n}$ yields a subset of $E^{\omega}$ and is defined by:
(Def. 2) There exists a function $c_{1}$ from $\mathbb{N}$ into $2^{E^{\omega}}$ such that $A^{n}=c_{1}(n)$ and $c_{1}(0)=\left\{\langle \rangle_{E}\right\}$ and for every $i$ holds $c_{1}(i+1)=c_{1}(i) \frown A$.

Next we state a number of propositions:
(23) $A^{n+1}=\left(A^{n}\right) \frown A$.
(24) $\quad A^{0}=\left\{\langle \rangle_{E}\right\}$.
(25) $A^{1}=A$.
(26) $A^{2}=A \frown A$.
(27) If $n \geq 1$, then $\left(\emptyset_{E^{\omega}}\right)^{n}=\emptyset$.
(28) $\left\{\left\rangle_{E}\right\}^{n}=\left\{\langle \rangle_{E}\right\}\right.$.
(29) $A^{n}=\left\{\langle \rangle_{E}\right\}$ iff $n=0$ or $A=\left\{\langle \rangle_{E}\right\}$.
(30) If $\left\rangle_{E} \in A\right.$, then $\left\rangle_{E} \in A^{n}\right.$.
(31) $\left(A^{n}\right) \frown A=A \frown A^{n}$.
(32) $A^{m+n}=\left(A^{m}\right) \frown A^{n}$.
(33) $\left(A^{m}\right)^{n}=A^{m \cdot n}$.
(34) If $\left\rangle_{E} \in A\right.$ and $n>0$, then $A \subseteq A^{n}$.
(35) If $\left\rangle_{E} \in A\right.$ and $n>0$ and $m>n$, then $A^{n} \subseteq A^{m}$.
(36) If $A \subseteq B$, then $A^{n} \subseteq B^{n}$.
(37) $A^{n} \cup B^{n} \subseteq(A \cup B)^{n}$.
(38) $(A \cap B)^{n} \subseteq A^{n} \cap B^{n}$.
(39) If $a \in C^{m}$ and $b \in C^{n}$, then $a{ }^{\wedge} b \in C^{m+n}$.

## 4. Closure of a Language

Let us consider $E, A$. The functor $A^{*}$ yielding a subset of $E^{\omega}$ is defined as follows:
(Def. 3) $\quad A^{*}=\bigcup\left\{B: \bigvee_{n} B=A^{n}\right\}$.
The following propositions are true:
(40) $x \in A^{*}$ iff there exists $n$ such that $x \in A^{n}$.
(41) $A^{n} \subseteq A^{*}$.
(42) If $x \in A$, then $x \in A^{*}$.
(43) $A \subseteq A^{*}$.
(44) $A \frown A \subseteq A^{*}$.
(45) If $a \in C^{*}$ and $b \in C^{*}$, then $a^{\curvearrowright} b \in C^{*}$.
(46) If $A \subseteq C^{*}$ and $B \subseteq C^{*}$, then $A \frown B \subseteq C^{*}$.
(47) $A^{*}=\left\{\langle \rangle_{E}\right\}$ iff $A=\emptyset$ or $A=\left\{\langle \rangle_{E}\right\}$.
(48) $\left\rangle_{E} \in A^{*}\right.$.
(49) If $A^{*}=\{x\}$, then $x=\langle \rangle_{E}$.
(50) If $x \in A^{m+1}$, then $x \in\left(A^{*}\right) \frown A$ and $x \in A \frown A^{*}$.
(51) If $x \in\left(A^{*}\right) \frown A$ or $x \in A \frown A^{*}$, then $x \in A^{*}$.
(52) If $\left\rangle_{E} \in A\right.$, then $A^{*}=\left(A^{*}\right) \frown A$ and $A^{*}=A \frown A^{*}$.
(53) If $\left\rangle_{E} \in A\right.$, then $A^{*}=\left(A^{*}\right) \frown A^{n}$ and $A^{*}=\left(A^{n}\right) \frown A^{*}$.
(54) $\quad A^{*}=\left\{\langle \rangle_{E}\right\} \cup A \frown A^{*}$ and $A^{*}=\left\{\langle \rangle_{E}\right\} \cup\left(A^{*}\right) \frown A$.
(55) $A \frown A^{*}=\left(A^{*}\right) \frown A$.
(56) $\left(A^{n}\right) \frown A^{*}=\left(A^{*}\right) \frown A^{n}$.
(57) If $A \subseteq B^{*}$, then $A^{n} \subseteq B^{*}$.
(58) If $A \subseteq B^{*}$, then $A^{*} \subseteq B^{*}$.
(59) If $A \subseteq B$, then $A^{*} \subseteq B^{*}$.
(60) $\left(A^{*}\right)^{*}=A^{*}$.
(61) $\left(A^{*}\right) \frown A^{*}=A^{*}$.
(62) $\quad\left(A^{n}\right)^{*} \subseteq A^{*}$.
(63) $\quad\left(A^{*}\right)^{n} \subseteq A^{*}$.
(64) If $n>0$, then $\left(A^{*}\right)^{n}=A^{*}$.
(65) If $A \subseteq B^{*}$, then $B^{*}=(B \cup A)^{*}$.
(66) If $a \in A^{*}$, then $A^{*}=(A \cup\{a\})^{*}$.
(67) $A^{*}=\left(A \backslash\left\{\langle \rangle_{E}\right\}\right)^{*}$.
(68) $\quad A^{*} \cup B^{*} \subseteq(A \cup B)^{*}$.
(69) $(A \cap B)^{*} \subseteq A^{*} \cap B^{*}$.
(70) $\langle x\rangle \in A^{*}$ iff $\langle x\rangle \in A$.

## 5. Alphabet as a Language

Let us consider $E$. The functor Lex $E$ yielding a subset of $E^{\omega}$ is defined by: (Def. 4) $\quad x \in$ Lex $E$ iff there exists $e$ such that $e \in E$ and $x=\langle e\rangle$.

Next we state three propositions:
(71) $a \in(\operatorname{Lex} E)^{\operatorname{len} a}$.
(72) $(\operatorname{Lex} E)^{*}=E^{\omega}$.
(73) If $A^{*}=E^{\omega}$, then Lex $E \subseteq A$.

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# Basic Properties of Determinants of Square Matrices over a Field ${ }^{1}$ 

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Summary. In this paper I present basic properties of the determinant of square matrices over a field and selected properties of the sign of a permutation. First, I define the sign of a permutation by the requirement

$$
\operatorname{sgn}(p)=\prod_{1 \leq i<j \leq n} \operatorname{sgn}(p(j)-p(i))
$$

where $p$ is any fixed permutation of a set with $n$ elements. I prove that the sign of a product of two permutations is the same as the product of their signs and show the relation between signs and parity of permutations. Then I consider the determinant of a linear combination of lines, the determinant of a matrix with permutated lines and the determinant of a matrix with a repeated line. Finally, at the end I prove that the determinant of a product of two square matrices is equal to the product of their determinants.

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The articles [21], [12], [27], [18], [13], [28], [7], [10], [8], [3], [4], [19], [25], [24], [16], [20], [11], [6], [5], [14], [22], [15], [31], [23], [26], [32], [1], [29], [9], [2], [17], and [30] provide the terminology and notation for this paper.

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## 1. The Sign of a Permutation

For simplicity, we use the following convention: $x, X$ denote sets, $i, j, k, l$, $n, m$ denote natural numbers, $D$ denotes a non empty set, $K$ denotes a field, $a, b$ denote elements of $K, p_{1}, p, q$ denote elements of the permutations of $n$-element set, $P_{1}, P$ denote permutations of $\operatorname{Seg} n, F$ denotes a function from $\operatorname{Seg} n$ into Seg $n, p_{2}, p_{3}, q_{2}, p_{4}$ denote elements of the permutations of $(n+2)$-element set, and $P_{2}$ denotes a permutation of $\operatorname{Seg}(n+2)$.

Let $X$ be a set. We introduce 2 Set $X$ as a synonym of TwoElementSets $(X)$.
The following three propositions are true:
(1) $\quad X \in 2 \operatorname{Set} \operatorname{Seg} n$ iff there exist $i, j$ such that $i \in \operatorname{Seg} n$ and $j \in \operatorname{Seg} n$ and $i<j$ and $X=\{i, j\}$.
(2) $2 \operatorname{Set} \operatorname{Seg} 0=\emptyset$ and $2 \operatorname{Set} \operatorname{Seg} 1=\emptyset$.
(3) For every $n$ such that $n \geq 2$ holds $\{1,2\} \in 2 \operatorname{Set} \operatorname{Seg} n$.

Let us consider $n$. Observe that $2 \operatorname{Set} \operatorname{Seg}(n+2)$ is non empty and finite.
Let us consider $n, x$ and let $p_{1}$ be an element of the permutations of $n$ element set. Note that $p_{1}(x)$ is natural.

Let us consider $K$. One can verify that the multiplication of $K$ is unital and the multiplication of $K$ is associative.

Let us consider $n, K$ and let $p_{2}$ be an element of the permutations of $(n+2)$ element set. The functor Part-sgn $\left(p_{2}, K\right)$ yielding a function from $2 \operatorname{Set} \operatorname{Seg}(n+2)$ into the carrier of $K$ is defined by the condition (Def. 1 ).
(Def. 1) Let $i, j$ be elements of $\mathbb{N}$ such that $i \in \operatorname{Seg}(n+2)$ and $j \in \operatorname{Seg}(n+2)$ and $i<j$. Then
(i) if $p_{2}(i)<p_{2}(j)$, then $\left(\operatorname{Part-sgn}\left(p_{2}, K\right)\right)(\{i, j\})=\mathbf{1}_{K}$, and
(ii) if $p_{2}(i)>p_{2}(j)$, then $\left(\operatorname{Part-sgn}\left(p_{2}, K\right)\right)(\{i, j\})=-\mathbf{1}_{K}$.

One can prove the following proposition
(4) Let $X$ be an element of Fin $2 \operatorname{Set} \operatorname{Seg}(n+2)$. Suppose that for every $x$ such that $x \in X$ holds $\left(\operatorname{Part-\operatorname {sgn}}\left(p_{3}, K\right)\right)(x)=\mathbf{1}_{K}$. Then (the multiplication of $K)-\sum_{X} \operatorname{Part}-\operatorname{sgn}\left(p_{3}, K\right)=\mathbf{1}_{K}$.
In the sequel $s$ denotes an element of $2 \operatorname{Set} \operatorname{Seg}(n+2)$.
The following propositions are true:
(5) $\quad\left(\operatorname{Part-sgn}\left(p_{3}, K\right)\right)(s)=\mathbf{1}_{K}$ or $\left(\operatorname{Part-sgn}\left(p_{3}, K\right)\right)(s)=-\mathbf{1}_{K}$.
(6) For all $i, j$ such that $i \in \operatorname{Seg}(n+2)$ and $j \in \operatorname{Seg}(n+2)$ and $i<j$ and $p_{3}(i)=q_{2}(i)$ and $p_{3}(j)=q_{2}(j)$ holds $\left(\operatorname{Part-sgn}\left(p_{3}, K\right)\right)(\{i, j\})=$ (Part-sgn $\left.\left(q_{2}, K\right)\right)(\{i, j\})$.
(7) Let $X$ be an element of $\operatorname{Fin} 2 \operatorname{Set} \operatorname{Seg}(n+2)$, given $p_{3}, q_{2}$, and $F$ be a finite set such that $F=\left\{s: s \in X \wedge\left(\operatorname{Part-sgn}\left(p_{3}, K\right)\right)(s) \neq\right.$ $\left.\left(\operatorname{Part-sgn}\left(q_{2}, K\right)\right)(s)\right\}$. Then
(i) if card $F \bmod 2=0$, then (the multiplication of $K)-\sum_{X} \operatorname{Part}-\operatorname{sgn}\left(p_{3}, K\right)=$ (the multiplication of $K)-\sum_{X} \operatorname{Part-\operatorname {sgn}}\left(q_{2}, K\right)$, and
(ii) if card $F \bmod 2=1$, then (the multiplication of $K)-\sum_{X} \operatorname{Part}-\operatorname{sgn}\left(p_{3}, K\right)=$ $-\left((\right.$ the multiplication of $\left.K)-\sum_{X} \operatorname{Part-sgn}\left(q_{2}, K\right)\right)$.
(8) Let $P$ be a permutation of $\operatorname{Seg} n$. Suppose $P$ is a transposition. Let given $i, j$. Suppose $i<j$. Then $P(i)=j$ if and only if the following conditions are satisfied:
(i) $i \in \operatorname{dom} P$,
(ii) $j \in \operatorname{dom} P$,
(iii) $P(i)=j$,
(iv) $P(j)=i$, and
(v) for every $k$ such that $k \neq i$ and $k \neq j$ and $k \in \operatorname{dom} P$ holds $P(k)=k$.
(9) Let given $p_{3}, q_{2}, p_{4}, i, j$. Suppose $p_{4}=p_{3} \cdot q_{2}$ and $q_{2}$ is a transposition and $q_{2}(i)=j$ and $i<j$. Let given $s$. If $\left(\operatorname{Part-sgn}\left(p_{3}, K\right)\right)(s) \neq$ $\left(\operatorname{Part}-\operatorname{sgn}\left(p_{4}, K\right)\right)(s)$, then $i \in s$ or $j \in s$.
(10) Let given $p_{3}, q_{2}, p_{4}, i, j, K$. Suppose $p_{4}=p_{3} \cdot q_{2}$ and $q_{2}$ is a transposition and $q_{2}(i)=j$ and $i<j$ and $\mathbf{1}_{K} \neq-\mathbf{1}_{K}$. Then
(i) $\quad\left(\operatorname{Part}-\operatorname{sgn}\left(p_{3}, K\right)\right)(\{i, j\}) \neq\left(\operatorname{Part}-\operatorname{sgn}\left(p_{4}, K\right)\right)(\{i, j\})$, and
(ii) for every $k$ such that $k \in \operatorname{Seg}(n+2)$ and $i \neq k$ and $j \neq k$ holds $\left(\operatorname{Part-sgn}\left(p_{3}, K\right)\right)(\{i, k\}) \neq\left(\operatorname{Part-sgn}\left(p_{4}, K\right)\right)(\{i, k\})$ iff $\left(\operatorname{Part-sgn}\left(p_{3}, K\right)\right)(\{j, k\}) \neq\left(\operatorname{Part-sgn}\left(p_{4}, K\right)\right)(\{j, k\})$.
Let us consider $n, K$ and let $p_{2}$ be an element of the permutations of $(n+2)$ element set. The functor $\operatorname{sgn}\left(p_{2}, K\right)$ yielding an element of $K$ is defined by:

The following propositions are true:
(11) $\operatorname{sgn}\left(p_{3}, K\right)=\mathbf{1}_{K}$ or $\operatorname{sgn}\left(p_{3}, K\right)=-\mathbf{1}_{K}$.
(12) For every element $I_{1}$ of the permutations of $(n+2)$-element set such that $I_{1}=\operatorname{idseq}(n+2)$ holds $\operatorname{sgn}\left(I_{1}, K\right)=\mathbf{1}_{K}$.
(13) For all $p_{3}, q_{2}, p_{4}$ such that $p_{4}=p_{3} \cdot q_{2}$ and $q_{2}$ is a transposition holds $\operatorname{sgn}\left(p_{4}, K\right)=-\operatorname{sgn}\left(p_{3}, K\right)$.
(14) For every element $t_{1}$ of the permutations of $(n+2)$-element set such that $t_{1}$ is a transposition holds $\operatorname{sgn}\left(t_{1}, K\right)=-\mathbf{1}_{K}$.
(15) Let $P$ be a finite sequence of elements of $A_{n+2}$ and $p_{3}$ be an element of the permutations of $(n+2)$-element set such that $p_{3}=\Pi P$ and for every $i$ such that $i \in \operatorname{dom} P$ there exists an element $t_{2}$ of the permutations of $(n+2)$-element set such that $P(i)=t_{2}$ and $t_{2}$ is a transposition. Then
(i) if len $P \bmod 2=0$, then $\operatorname{sgn}\left(p_{3}, K\right)=\mathbf{1}_{K}$, and
(ii) if len $P \bmod 2=1$, then $\operatorname{sgn}\left(p_{3}, K\right)=-\mathbf{1}_{K}$.
(16) Let given $i, j, n$. Suppose $i<j$ and $i \in \operatorname{Seg} n$ and $j \in \operatorname{Seg} n$. Then there exists an element $t_{1}$ of the permutations of $n$-element set such that $t_{1}$ is a
transposition and $t_{1}(i)=j$.
(17) Let $p$ be an element of the permutations of $(k+1)$-element set. Suppose $p(k+1) \neq k+1$. Then there exists an element $t_{1}$ of the permutations of $(k+1)$-element set such that $t_{1}$ is a transposition and $t_{1}(p(k+1))=k+1$ and $\left(t_{1} \cdot p\right)(k+1)=k+1$.
(18) Let given $X, x$. Suppose $x \notin X$. Let $p_{5}$ be a permutation of $X \cup\{x\}$. If $p_{5}(x)=x$, then there exists a permutation $p$ of $X$ such that $p_{5} \upharpoonright X=p$.
(19) Let $p, q$ be permutations of $X$ and $p_{5}, q_{1}$ be permutations of $X \cup\{x\}$. If $p_{5} \upharpoonright X=p$ and $q_{1} \upharpoonright X=q$ and $p_{5}(x)=x$ and $q_{1}(x)=x$, then $\left(p_{5} \cdot q_{1}\right) \upharpoonright X=$ $p \cdot q$ and $\left(p_{5} \cdot q_{1}\right)(x)=x$.
(20) For every element $t_{1}$ of the permutations of $k$-element set such that $t_{1}$ is a transposition holds $t_{1} \cdot t_{1}=\operatorname{idseq}(k)$ and $t_{1}=t_{1}{ }^{-1}$.
(21) Let given $p_{1}$. Then there exists a finite sequence $P$ of elements of $A_{n}$ such that
(i) $p_{1}=\prod P$, and
(ii) for every $i$ such that $i \in \operatorname{dom} P$ there exists an element $t_{2}$ of the permutations of $n$-element set such that $P(i)=t_{2}$ and $t_{2}$ is a transposition.
(22) $K$ is Fanoian iff $\mathbf{1}_{K} \neq-\mathbf{1}_{K}$.
(23) For every Fanoian field $K$ holds $p_{2}$ is even iff $\operatorname{sgn}\left(p_{2}, K\right)=\mathbf{1}_{K}$ and $p_{2}$ is odd iff $\operatorname{sgn}\left(p_{2}, K\right)=-\mathbf{1}_{K}$.
(24) For all $p_{3}, q_{2}, p_{4}$ such that $p_{4}=p_{3} \cdot q_{2}$ holds $\operatorname{sgn}\left(p_{4}, K\right)=\operatorname{sgn}\left(p_{3}, K\right)$. $\operatorname{sgn}\left(q_{2}, K\right)$.
(25) $p$ is even and $q$ is even or $p$ is odd and $q$ is odd iff $p \cdot q$ is even.
(26) $(-1)^{\operatorname{sgn}\left(p_{2}\right)} a=\operatorname{sgn}\left(p_{2}, K\right) \cdot a$.
(27) For every element $t_{1}$ of the permutations of $(n+2)$-element set such that $t_{1}$ is a transposition holds $t_{1}$ is odd.

Let us consider $n$. Observe that there exists a permutation of $\operatorname{Seg}(n+2)$ which is odd.

## 2. The Determinant of a Linear Combination of Lines

For simplicity, we follow the rules: $p_{6}$ denotes a finite sequence of elements of $D, M$ denotes a matrix over $D$ of dimension $n \times m, p_{7}, q_{3}$ denote finite sequences of elements of $K$, and $A, B$ denote matrices over $K$ of dimension $n$.

Let us consider $l, n, m, D$, let $M$ be a matrix over $D$ of dimension $n \times m$, and let $p_{6}$ be a finite sequence of elements of $D$. The functor $\operatorname{ReplaceLine}\left(M, l, p_{6}\right)$ yields a matrix over $D$ of dimension $n \times m$ and is defined as follows:
(Def. 3)(i) len ReplaceLine $\left(M, l, p_{6}\right)=\operatorname{len} M$ and width $\operatorname{ReplaceLine}\left(M, l, p_{6}\right)=$ width $M$ and for all $i, j$ such that $\langle i, j\rangle \in$ the indices of $M$ holds
if $i \neq l$, then $\left(\operatorname{ReplaceLine}\left(M, l, p_{6}\right)\right)_{i, j}=M_{i, j}$ and if $i=l$, then (ReplaceLine $\left.\left(M, l, p_{6}\right)\right)_{l, j}=p_{6}(j)$ if len $p_{6}=\operatorname{width} M$,
(ii) $\operatorname{ReplaceLine}\left(M, l, p_{6}\right)=M$, otherwise.

Let us consider $l, n, m, D$, let $M$ be a matrix over $D$ of dimension $n \times m$, and let $p_{6}$ be a finite sequence of elements of $D$. We introduce $\operatorname{RLine}\left(M, l, p_{6}\right)$ as a synonym of ReplaceLine $\left(M, l, p_{6}\right)$.

The following propositions are true:
(28) For all $l, M, p_{6}, i$ such that $i \in \operatorname{Seg} n$ holds if $i=l$ and len $p_{6}=\operatorname{width} M$, then $\operatorname{Line}\left(\operatorname{RLine}\left(M, l, p_{6}\right), i\right)=p_{6}$ and if $i \neq l$, then $\operatorname{Line}\left(\operatorname{RLine}\left(M, l, p_{6}\right), i\right)=\operatorname{Line}(M, i)$.
(29) For all $M, p_{6}$ such that len $p_{6}=$ width $M$ and for every element $p^{\prime}$ of $D^{*}$ such that $p_{6}=p^{\prime}$ holds $\operatorname{RLine}\left(M, l, p_{6}\right)=\operatorname{Replace}\left(M, l, p^{\prime}\right)$.
(30) $\quad M=\operatorname{RLine}(M, l, \operatorname{Line}(M, l))$.
(31) Let given $l, p_{7}, q_{3}, p_{1}$. Suppose $l \in \operatorname{Seg} n$ and $\operatorname{len} p_{7}=n$ and len $q_{3}=n$. Let $M$ be a matrix over $K$ of dimension $n$. Then (the multiplication of $K) \circledast\left(p_{1}-\operatorname{Path} \operatorname{RLine}\left(M, l, a \cdot p_{7}+b \cdot q_{3}\right)\right)=a \cdot(($ the multiplication of $\left.K) \circledast\left(p_{1}-\operatorname{Path} \operatorname{RLine}\left(M, l, p_{7}\right)\right)\right)+b \cdot(($ the multiplication of $\left.K) \circledast\left(p_{1}-\operatorname{Path} \operatorname{RLine}\left(M, l, q_{3}\right)\right)\right)$.
(32) Let given $l, p_{7}, q_{3}, p_{1}$. Suppose $l \in \operatorname{Seg} n$ and len $p_{7}=n$ and len $q_{3}=$ $n$. Let $M$ be a matrix over $K$ of dimension $n$. Then (the product on paths of $\left.\operatorname{RLine}\left(M, l, a \cdot p_{7}+b \cdot q_{3}\right)\right)\left(p_{1}\right)=a \cdot$ (the product on paths of $\left.\operatorname{RLine}\left(M, l, p_{7}\right)\right)\left(p_{1}\right)+b \cdot\left(\right.$ the product on paths of $\left.\operatorname{RLine}\left(M, l, q_{3}\right)\right)\left(p_{1}\right)$.
(33) Let given $l, p_{7}, q_{3}$. Suppose $l \in \operatorname{Seg} n$ and len $p_{7}=n$ and len $q_{3}=n$. Let $M$ be a matrix over $K$ of dimension $n$. Then $\operatorname{Det} \operatorname{RLine}\left(M, l, a \cdot p_{7}+b \cdot q_{3}\right)=$ $a \cdot \operatorname{Det} \operatorname{RLine}\left(M, l, p_{7}\right)+b \cdot \operatorname{Det} \operatorname{RLine}\left(M, l, q_{3}\right)$.
(34) If $l \in \operatorname{Seg} n$ and $\operatorname{len} p_{7}=n$, then $\operatorname{Det} \operatorname{RLine}\left(A, l, a \cdot p_{7}\right)=a$. $\operatorname{Det} \operatorname{RLine}\left(A, l, p_{7}\right)$.
(35) If $l \in \operatorname{Seg} n$, then $\operatorname{Det} \operatorname{RLine}(A, l, a \cdot \operatorname{Line}(A, l))=a \cdot \operatorname{Det} A$.
(36) If $l \in \operatorname{Seg} n$ and len $p_{7}=n$ and len $q_{3}=n$, then $\operatorname{Det} \operatorname{RLine}\left(A, l, p_{7}+q_{3}\right)=$ $\operatorname{Det} \operatorname{RLine}\left(A, l, p_{7}\right)+\operatorname{Det} \operatorname{RLine}\left(A, l, q_{3}\right)$.

## 3. The Determinant of a Matrix with Permutated Lines and with a Repeated Line

Let us consider $n, m, D$, let $F$ be a function from $\operatorname{Seg} n \operatorname{into} \operatorname{Seg} n$, and let $M$ be a matrix over $D$ of dimension $n \times m$. Then $M \cdot F$ is a matrix over $D$ of dimension $n \times m$ and it can be characterized by the condition:
(Def. 4) $\operatorname{len}(M \cdot F)=\operatorname{len} M$ and $\operatorname{width}(M \cdot F)=\operatorname{width} M$ and for all $i, j, k$ such that $\langle i, j\rangle \in$ the indices of $M$ and $F(i)=k$ holds $(M \cdot F)_{i, j}=M_{k, j}$.
The following propositions are true:
(37)(i) The indices of $M=$ the indices of $M \cdot F$, and
(ii) for all $i, j$ such that $\langle i, j\rangle \in$ the indices of $M$ there exists $k$ such that $F(i)=k$ and $\langle k, j\rangle \in$ the indices of $M$ and $(M \cdot F)_{i, j}=M_{k, j}$.
(38) For every matrix $M$ over $D$ of dimension $n \times m$ and for every $F$ and for every $k$ such that $k \in \operatorname{Seg} n$ holds $\operatorname{Line}(M \cdot F, k)=M(F(k))$.
(39) $M \cdot \operatorname{idseq}(n)=M$.
(40) For all $p, P_{1}, q$ such that $q=p \cdot P_{1}^{-1}$ holds $p$-Path $A \cdot P_{1}=(q-\operatorname{Path} A) \cdot P_{1}$.
(41) For all $p, P_{1}, q$ such that $q=p \cdot P_{1}^{-1}$ holds (the multiplication of $K) \circledast\left(p\right.$-Path $\left.A \cdot P_{1}\right)=($ the multiplication of $K) \circledast(q$-Path $A)$.
(42) For all $p_{3}, q_{2}$ such that $q_{2}=p_{3}^{-1} \operatorname{holds} \operatorname{sgn}\left(p_{3}, K\right)=\operatorname{sgn}\left(q_{2}, K\right)$.
(43) Let $M$ be a matrix over $K$ of dimension $n+2$ and given $p_{2}, P_{2}$. Suppose $p_{2}=P_{2}$. Let given $p_{3}, q_{2}$. Suppose $q_{2}=p_{3} \cdot P_{2}{ }^{-1}$. Then (the product on paths of $M)\left(q_{2}\right)=\operatorname{sgn}\left(p_{2}, K\right) \cdot\left(\right.$ the product on paths of $\left.M \cdot P_{2}\right)\left(p_{3}\right)$.
(44) Let given $p_{1}$. Then there exists a permutation $P$ of the permutations of $n$-element set such that for every element $p$ of the permutations of $n$-element set holds $P(p)=p \cdot p_{1}$.
(45) For every matrix $M$ over $K$ of dimension $n+2 \times n+2$ and for all $p_{2}$, $P_{2}$ such that $p_{2}=P_{2}$ holds $\operatorname{Det}\left(M \cdot P_{2}\right)=\operatorname{sgn}\left(p_{2}, K\right) \cdot \operatorname{Det} M$.
(46) For every matrix $M$ over $K$ of dimension $n$ and for all $p_{1}, P_{1}$ such that $p_{1}=P_{1}$ holds $\operatorname{Det}\left(M \cdot P_{1}\right)=(-1)^{\operatorname{sgn}\left(p_{1}\right)} \operatorname{Det} M$.
(47) Let $P_{3}$ be a permutation of the permutations of $n$-element set and given $p_{1}$. If $p_{1}$ is odd and for every $p$ holds $P_{3}(p)=p \cdot p_{1}$, then $P_{3}{ }^{\circ}\{p: p$ is even $\}=\{q: q$ is odd $\}$.
(48) Let given $n$. Suppose $n \geq 2$. Then there exist finite sets $O_{1}, E_{1}$ such that $E_{1}=\{p: p$ is even $\}$ and $O_{1}=\{q: q$ is odd $\}$ and $E_{1} \cap O_{1}=\emptyset$ and $E_{1} \cup O_{1}=$ the permutations of $n$-element set and card $E_{1}=\operatorname{card} O_{1}$.
(49) Let given $i, j$. Suppose $i \in \operatorname{Seg} n$ and $j \in \operatorname{Seg} n$ and $i<j$. Let $M$ be a matrix over $K$ of dimension $n$. Suppose $\operatorname{Line}(M, i)=\operatorname{Line}(M, j)$. Let $p$, $q, t_{1}$ be elements of the permutations of $n$-element set. Suppose $q=p \cdot t_{1}$ and $t_{1}$ is a transposition and $t_{1}(i)=j$. Then (the product on paths of $M)(q)=-($ the product on paths of $M)(p)$.
(50) Let given $i, j$. Suppose $i \in \operatorname{Seg} n$ and $j \in \operatorname{Seg} n$ and $i<j$. Let $M$ be a matrix over $K$ of dimension $n$. If $\operatorname{Line}(M, i)=\operatorname{Line}(M, j)$, then Det $M=0_{K}$.
(51) For all $i, j$ such that $i \in \operatorname{Seg} n$ and $j \in \operatorname{Seg} n$ and $i \neq j$ holds $\operatorname{Det} \operatorname{RLine}(A, i, \operatorname{Line}(A, j))=0_{K}$.
(52) For all $i, j$ such that $i \in \operatorname{Seg} n$ and $j \in \operatorname{Seg} n$ and $i \neq j$ holds $\operatorname{Det} \operatorname{RLine}(A, i, a \cdot \operatorname{Line}(A, j))=0_{K}$.
(53) For all $i, j$ such that $i \in \operatorname{Seg} n$ and $j \in \operatorname{Seg} n$ and $i \neq j$ holds $\operatorname{Det} A=$
$\operatorname{Det} \operatorname{RLine}(A, i, \operatorname{Line}(A, i)+a \cdot \operatorname{Line}(A, j))$.
If $F \notin$ the permutations of $n$-element set, then $\operatorname{Det}(A \cdot F)=0_{K}$.

## 4. The Determinant of a Product of Two Square Matrices

Let $K$ be a non empty loop structure. The functor addFinS $K$ yielding a binary operation on (the carrier of $K)^{*}$ is defined as follows:
(Def. 5) For all elements $p_{5}, p_{3}$ of (the carrier of $\left.K\right)^{*}$ holds (addFinS $\left.K\right)\left(p_{5}\right.$, $\left.p_{3}\right)=p_{5}+p_{3}$.
Let $K$ be an Abelian non empty loop structure. One can verify that addFinS $K$ is commutative.

Let $K$ be an add-associative non empty loop structure. Note that addFinS $K$ is associative.

The following propositions are true:
(55) Let $A, B$ be matrices over $K$. Suppose width $A=\operatorname{len} B$ and len $B>0$. Let given $i$. Suppose $i \in \operatorname{Seg}$ len $A$. Then there exists a finite sequence $P$ of elements of (the carrier of $K)^{*}$ such that len $P=\operatorname{len} B$ and $\operatorname{Line}(A$. $B, i)=\operatorname{addFinS} K \odot P$ and for every $j$ such that $j \in \operatorname{Seg}$ len $B$ holds $P(j)=A_{i, j} \cdot \operatorname{Line}(B, j)$.
(56) Let $A, B, C$ be matrices over $K$ of dimension $n$ and given $i$. Suppose $i \in \operatorname{Seg} n$. Then there exists a finite sequence $P$ of elements of $K$ such that len $P=n$ and $\operatorname{Det} \operatorname{RLine}(C, i, \operatorname{Line}(A \cdot B, i))=$ the addition of $K \odot P$ and for every $j$ such that $j \in \operatorname{Seg} n$ holds $P(j)=$ $A_{i, j} \cdot \operatorname{Det} \operatorname{RLine}(C, i, \operatorname{Line}(B, j))$.
(57) Let $X$ be a set, $Y$ be a non empty set, and given $x$. Suppose $x \notin X$. Then there exists a function $B_{1}$ from : : $Y^{X}, Y$ : into $Y^{X \cup\{x\}}$ such that
(i) $\quad B_{1}$ is bijective, and
(ii) for every function $f$ from $X$ into $Y$ and for every function $F$ from $X \cup\{x\}$ into $Y$ such that $F \upharpoonright X=f$ holds $B_{1}(\langle f, F(x)\rangle)=F$.
(58) Let $X$ be a finite set, $Y$ be a non empty finite set, and given $x$. Suppose $x \notin X$. Let $F$ be a binary operation on $D$. Suppose $F$ is commutative and associative and has a unity and an inverse operation. Let $f$ be a function from $Y^{X}$ into $D$ and $g$ be a function from $Y^{X \cup\{x\}}$ into $D$. Suppose that for every function $H$ from $X$ into $Y$ and for every element $S_{1}$ of $\operatorname{Fin}\left(Y^{X \cup\{x\}}\right)$ such that $S_{1}=\{h ; h$ ranges over functions from $X \cup\{x\}$ into $Y: h \upharpoonright X=H\}$ holds $F-\sum_{S_{1}} g=f(H)$. Then $F-\sum_{\Omega_{Y X}^{\mathrm{f}}} f=F-\sum_{\Omega_{Y}^{\mathrm{f}}} f \cup\{x\}$.
(59) Let $A, B$ be matrices over $D$ of dimension $n \times m$ and given $i$. Suppose $i \leq n$ and $0<n$. Let $F$ be a function from $\operatorname{Seg} i$ into $\operatorname{Seg} n$. Then there exists a matrix $M$ over $D$ of dimension $n \times m$ such that $M=A+\cdot(B$.
(idseq $(n)+\cdot F)) \upharpoonright \operatorname{Seg} i$ and for every $j$ holds if $j \in \operatorname{Seg} i$, then $M(j)=$ $B(F(j))$ and if $j \notin \operatorname{Seg} i$, then $M(j)=A(j)$.
(60) Let $A, B$ be matrices over $K$ of dimension $n$. Suppose $0<n$. Then there exists a function $P$ from $(\operatorname{Seg} n)^{\operatorname{Seg} n}$ into the carrier of $K$ such that
(i) for every function $F$ from $\operatorname{Seg} n$ into $\operatorname{Seg} n$ there exists a finite sequence $P_{4}$ of elements of $K$ such that len $P_{4}=n$ and for all natural numbers $F_{1}, j$ such that $j \in \operatorname{Seg} n$ and $F_{1}=F(j)$ holds $P_{4}(j)=A_{j, F_{1}}$ and $P(F)=(($ the multiplication of $\left.K) \circledast\left(P_{4}\right)\right) \cdot \operatorname{Det}(B \cdot F)$, and
(ii) $\operatorname{Det}(A \cdot B)=($ the addition of $K)-\sum_{\Omega_{(\operatorname{Seg} n)}^{\mathrm{f} \operatorname{Seg} n}} P$.
(61) Let $A, B$ be matrices over $K$ of dimension $n$. Suppose $0<n$. Then there exists a function $P$ from the permutations of $n$-element set into the carrier of $K$ such that
(i) $\operatorname{Det}(A \cdot B)=($ the addition of $K)-\sum_{\Omega_{\text {the permutations of } n \text {-element set }}^{\mathrm{f}}} P$, and
(ii) for every element $p_{1}$ of the permutations of $n$-element set holds $P\left(p_{1}\right)=$ $\left((\right.$ the multiplication of $K) \circledast\left(p_{1}\right.$-Path $\left.\left.A\right)\right) \cdot(-1)^{\operatorname{sgn}\left(p_{1}\right)} \operatorname{Det} B$.
(62) For all matrices $A, B$ over $K$ of dimension $n$ such that $0<n$ holds $\operatorname{Det}(A \cdot B)=\operatorname{Det} A \cdot \operatorname{Det} B$.

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